

# [MS-PATCH]: LZX DELTA Compression and Decompression

## Intellectual Property Rights Notice for Protocol Documentation

- **Copyrights.** This protocol documentation is covered by Microsoft copyrights. Regardless of any other terms that are contained in the terms of use for the Microsoft website that hosts this documentation, you may make copies of it in order to develop implementations of the protocols, and may distribute portions of it in your implementations of the protocols or your documentation as necessary to properly document the implementation. This permission also applies to any documents that are referenced in the protocol documentation.
- **No Trade Secrets.** Microsoft does not claim any trade secret rights in this documentation.
- **Patents.** Microsoft has patents that may cover your implementations of the protocols. Neither this notice nor Microsoft's delivery of the documentation grants any licenses under those or any other Microsoft patents. However, the protocols may be covered by Microsoft's Open Specification Promise (available here: <http://www.microsoft.com/interop/osp/default.aspx>). If you would prefer a written license, or if the protocols are not covered by the OSP, patent licenses are available by contacting [protocol@microsoft.com](mailto:protocol@microsoft.com).
- **Trademarks.** The names of companies and products contained in this documentation may be covered by trademarks or similar intellectual property rights. This notice does not grant any licenses under those rights.

**Reservation of Rights.** All other rights are reserved, and this notice does not grant any rights other than specifically described above, whether by implication, estoppel, or otherwise.

**Preliminary Documentation.** This documentation is preliminary documentation for these protocols. Since the documentation may change between this preliminary version and the final version, there are risks in relying on preliminary documentation. To the extent that you incur additional development obligations or any other costs as a result of relying on this preliminary documentation, you do so at your own risk.

**Tools.** This protocol documentation is intended for use in conjunction with publicly available standard specifications and networking programming art, and assumes that the reader is either familiar with the aforementioned material or has immediate access to it. A protocol specification does not require the use of Microsoft programming tools or programming environments in order for a Licensee to develop an implementation. Licensees who have access to Microsoft programming tools and environments are free to take advantage of them.

Revision Summary			
Author	Date	Version	Comments
Microsoft Corporation	April 4, 2008	0.1	Initial Availability

Preliminary

## Table of Contents

<b>1</b>	<b>Introduction</b> .....	<b>4</b>
1.1	Glossary .....	4
1.2	References.....	4
1.2.1	Normative References .....	4
1.2.2	Informative References .....	4
<b>2</b>	<b>Description</b> .....	<b>4</b>
2.1	LZ77.....	4
2.2	LZX.....	5
2.3	LZXD.....	5
2.4	Bitstream.....	5
2.5	Window Size .....	5
2.6	Reference Data .....	5
2.7	Huffman Trees.....	6
2.8	Position Slot.....	7
2.9	Repeated Offsets.....	7
2.10	Match Lengths .....	8
2.11	E8 Call Translation.....	8
2.12	Chunk Size.....	10
2.13	Block Header .....	11
2.14	Block Type .....	11
2.15	Block Size.....	11
2.15.1	Uncompressed Block .....	12
2.15.2	Verbatim Block .....	12
2.15.3	Aligned Offset Block .....	13
2.15.4	Encoding the Trees and Pre-Trees.....	13
2.15.5	Compressed Token Sequence.....	14
2.15.6	Converting Match Offset into Formatted Offset Values .....	15
2.15.7	Converting Formatted Offset into Position Slot and Position Footer Values....	16
2.15.8	Converting Position Footer into Verbatim Bits or Aligned Offset Bits .....	17
2.15.9	Converting Match Length into Length Header and Length Footer Values.....	18
2.15.10	Converting Length Header and Position Slot into Length/Position Header Values	19
2.16	Extra Length .....	19
2.16.1	Encoding a Match.....	19
2.17	Encoding a Literal .....	19
2.17.1	Decoding Matches and Literals (Aligned and Verbatim Blocks).....	20
<b>3</b>	<b>Protocol Examples</b> .....	<b>21</b>
<b>4</b>	<b>Appendix A: Office/Exchange Behavior</b> .....	<b>22</b>
	<b>Index</b> .....	<b>23</b>

# 1 Introduction

LZX is an LZ77-based Microsoft compression engine described in the Microsoft Cabinet SDK. LZXD (D for Delta) is a derivative of the Microsoft Cabinet LZX format with some modifications to facilitate efficient delta compression. Delta compression is a technique in which one set of data can be compressed within the context of a reference set of data that is supplied both to the compressor and decompressor. Delta compression is commonly used to encode updates to similar existing data sets so that the size of compressed data can be significantly reduced relative to ordinary non-delta compression techniques. Expanding a delta-compressed set of data requires that the exact same reference data be provided during decompression.

## 1.1 Glossary

The following terms are specific to this document:

**MAY, SHOULD, MUST, SHOULD NOT, MUST NOT:** These terms (in all caps) are used as described in [RFC2119]. All statements of optional behavior use either MAY, SHOULD, or SHOULD NOT.

## 1.2 References

### 1.2.1 Normative References

[MS-OXGLOS] Microsoft Corporation, "Office Exchange Protocols Master Glossary", April 2008.

[RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, March 1997, <http://www.ietf.org/rfc/rfc2119.txt>.

### 1.2.2 Informative References

None.

# 2 Description

## 2.1 LZ77

LZ77 refers to the well-known Lempel-Ziv 1977 sliding window data compression algorithm.

## 2.2 LZX

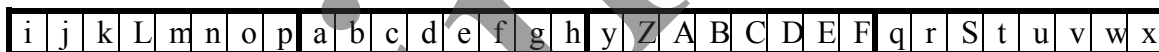
LZX is an LZ77-based compressor that uses static Huffman encoding and a sliding window of selectable size. LZX is most commonly known as part of the Microsoft Cabinet compression format. Data symbols are encoded either as an uncompressed symbol, or as a logical (offset, length) pair indicating that length symbols shall be copied from a displacement of offset symbols from the current position in the output stream. The value of offset is constrained to be less than the current position in the output stream, up to the size of the sliding window.

## 2.3 LZXD

LZXD is an LZX variant modified to facilitate efficient delta-compression. LZXD provides a mechanism for both compressor and decompressor to refer to a common reference set of data, and relaxes the constraint that match offset be constrained to less than the current position in the output stream, allowing match offset to refer to the logically prepended reference data. This effectively enables the compressed data stream to encode “matches” both from the reference data and from the uncompressed data stream.

## 2.4 Bitstream

An LZXD Bitstream is encoded as a sequence of aligned 16-bit integers stored in the order least-significant-byte most-significant-byte, also known as byte-swapped or little-endian words. Given an input stream of bits named a, b, c, ..., x, y, z, A, B, C, D, E, F, the output byte stream (with byte boundaries highlighted) would be as shown below.



## 2.5 Window Size

The sliding window size MUST be a power of 2, from  $2^{17}$  (128 KB) up to  $2^{25}$  (32 MB). The window size is not stored in the compressed data stream, and MUST be specified to the decoder before decoding begins. The preferred window size is the smallest power of two between  $2^{17}$  and  $2^{25}$  that is greater than or equal to the sum of the size of the reference data rounded up to multiple of 32,768 and the size of the subject data.

## 2.6 Reference Data

For delta compression, the reference data is a sequence of bytes given to the compressor prior to compressing the subject data. The exact same reference data sequence MUST be given to the decompressor prior to decompression. The reference data sequence is treated as logically prepended to the subject data sequence being compressed or decompressed. During decompression, match offsets are negative displacements from the “current position” in the

output stream, up to the specified Window Size. When match offset values exceed the number of bytes already emitted in the uncompressed output stream, they are simply pointing into the reference data that is logically prepended to the subject data.

Offset	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
Value	A	B	C	D	E	F	G	H	I	J	a	b	c	D	E	F	a	b	c	e
<b>Reference Data Sequence</b>											<b>Subject Data Sequence</b>									

In this example, the reference data is 10 bytes long and consists of the sequence “ABCDEFGHJI”. The data to be compressed, or the subject data, is also 10 bytes long (although the data does not need to be the same length as the reference data) and consists of “abcDEFabce”. A valid encoded sequence would consist of the following tokens:

‘a’, ‘b’, ‘c’, (match offset -10, length 3), (match offset -6, length 3), ‘e’

The first match offset exceeds the amount of subject data already in the window, pointing instead into the reference data portion. The second match offset does not exceed the amount of subject data in the window and instead refers to a portion of the subject data previously compressed or decompressed.

## 2.7 Huffman Trees

LZXD uses canonical Huffman tree structures to represent elements. Huffman trees are well known in data compression and are not described here. Because an LZXD decoder uses only the path lengths of the Huffman tree to reconstruct the identical tree, the following constraints are made on the tree structure.

For any two elements with the same path length, the lower-numbered element MUST be further left on the tree than the higher numbered element. An alternative way of stating this constraint is that lower-numbered elements MUST have lower path traversal values; for example, 0010 (left-left-right-left) is lower than 0011 (left-left-right-right).

For each level, starting at the deepest level of the tree and then moving upwards, leaf nodes MUST start as far left as possible. An alternative way of stating this constraint is that if any tree node has children then all tree nodes to the left of it with the same path length MUST also have children.

A non-empty Huffman tree MUST contain at least 2 elements, so in the case where all but one tree element has zero frequency, the resulting tree MUST minimally consist of two Huffman codes, “0” and “1”.

LZXD uses several Huffman tree structures. The Main Tree comprises 256 elements corresponding to all possible 8-bit characters, plus 8 \* NUM\_POSITION\_SLOTS elements corresponding to matches. The value of NUM\_POSITION\_SLOTS depends on the specified window size as described in the next section. The Length Tree comprises 249 elements. Other

6 of 24

trees, such as the Aligned Offset Tree (comprising 8 elements), and the Pre-Trees (comprising 20 elements each), have a smaller role.

## 2.8 Position Slot

The window size determines the number of window subdivisions, or “position slots”, as shown in the following table:

**Window Size/Position Slot Table**

Window size	Position slots required
128 KB	34
256 KB	36
512 KB	38
1 MB	42
2 MB	50
4 MB	66
8 MB	98
16 MB	162
32 MB	290

## 2.9 Repeated Offsets

LZXD extends the conventional LZ77 format in several ways, one of which is in the use of repeated offset codes. Three match offset codes, named the repeated offset codes, are reserved to indicate that the current match offset is the same as that of one of the three previous matches which is not itself a repeated offset.

The three special offset codes are encoded as offset values 0, 1, and 2 (for example, encoding an offset of 0 means “use the most recent non-repeated match offset”, an offset of 1 means “use the second most recent non-repeated match offset”, etc.). All remaining offset values are displaced by +3, as is shown in the table below, which prevents matches at offsets WINDOW\_SIZE, WINDOW\_SIZE-1, and WINDOW\_SIZE-2.

**Correlation Between Encoded Offset and Real Offset**

Encoded offset	Real offset
0	Most recent real match offset
1	Second most recent match offset
2	Third most recent match offset
3	1 (closest allowable)
4	2
5	3

6	4
7	5
8	6
500	498
x+2	X
WINDOW_SIZE-1 (maximum possible)	WINDOW_SIZE-3

The three most recent real match offsets are kept in a list, the behavior of which explained below:

Let R0 be defined as the most recent real offset

Let R1 be defined as the second most recent offset

Let R2 be defined as the third most recent offset

The list is managed similarly to an LRU (least recently used) queue, with the exception of the cases when R1 or R2 is output. In these cases, R1 or R2 is simply swapped with R0, which requires fewer operations than would an LRU queue.

The initial state of R0, R1, R2 is (1, 1, 1).

#### Management of the Repeated Offsets List

Match offset X where...	Operation
$X \neq R0$ and $X \neq R1$ and $X \neq R2$	$R2 \leftarrow R1$ $R1 \leftarrow R0$ $R0 \leftarrow X$
$X = R0$	None
$X = R1$	swap $R0 \leftrightarrow R1$
$X = R2$	swap $R0 \leftrightarrow R2$

## 2.10 Match Lengths

The minimum match length (number of bytes) encoded by LZXD is 2 bytes, and the maximum match length is 32,768 bytes. However, no match of any length can span a modulo-32 KB boundary in the uncompressed stream. Match length encoding is combined with match position encoding as described in the Compressed Token Sequence section below.

## 2.11 E8 Call Translation

E8 Call Translation is an optional feature that is sometimes used when the data to compress contains x86 instruction sequences. E8 Translation operates as a pre-processing stage prior to compressing each chunk, and the compressed stream header contains a bit that indicates



whether the decoder shall reverse the translation as a post-processing step after decompressing each chunk.

The x86 instruction beginning with a byte value of 0xE8 is followed by a 32-bit little-endian relative displacement to the call target. When E8 Call Translation is enabled, the following pre-processing step is performed on the uncompressed input prior to compression (assuming little-endian byte ordering):

Let `chunk_offset` refer to the total number of uncompressed bytes preceding this chunk.  
Let `E8_file_size` refer to the caller-specified value given to the compressor or decoded from the header of the compressed stream during decompression.

For each 32 KB chunk of uncompressed data (or less than 32 KB if last chunk to compress):

```
if ( ( chunk_offset < 0x40000000 ) && ( chunk_size > 10 ) )
  for ( i = 0; i < ( chunk_size - 10 ); i++ )
    if ( chunk_byte[ i ] == 0xE8 )
      long current_pointer = chunk_offset + i;
      long displacement =  chunk_byte[ i+1 ] |
                          chunk_byte[ i+2 ] << 8 |
                          chunk_byte[ i+3 ] << 16 |
                          chunk_byte[ i+4 ] << 24;

      long target = current_pointer + displacement;
      if ( ( target >= 0 ) && ( target <
E8_file_size+current_pointer) )
        if ( target >= E8_file_size )
          target = displacement - E8_file_size;
        endif
        chunk_byte[ i+1 ] = (byte)( target );
        chunk_byte[ i+2 ] = (byte)( target >> 8 );
        chunk_byte[ i+3 ] = (byte)( target >> 16 );
        chunk_byte[ i+4 ] = (byte)( target >> 24 );
      endif
      i += 4;
    endif
  endfor
endif
```

After decompression, the E8 scanning algorithm is the same, but the translation reversal is:

```

long value = chunk_byte[ i+1 ]      |
             chunk_byte[ i+2 ] << 8 |
             chunk_byte[ i+3 ] << 16 |
             chunk_byte[ i+4 ] << 24;

if ( ( value >= -current_pointer ) && ( value <
E8_file_size ) )
    if ( ( value >= 0 ) )
        displacement = value - current_pointer;
    else
        displacement = value + E8_file_size;
    endif
    chunk_byte[ i+1 ] = (byte)( displacement );
    chunk_byte[ i+2 ] = (byte)( displacement >> 8 );
    chunk_byte[ i+3 ] = (byte)( displacement >> 16 );
    chunk_byte[ i+4 ] = (byte)( displacement >> 24 );
endif

```

The first bit in the first Chunk in the LZXD bitstream (following the 2-byte Chunk Size prefix described below) indicates the presence or absence of two 16-bit fields immediately following the single bit. If the bit is set, E8 translation is enabled using the 32-bit value derived from the two 16-bit fields as the `E8_file_size` provided to the compressor when E8 translation was enabled. Note that `E8_file_size` is completely independent of the length of the uncompressed data. E8 call translation is always disabled after the 32,768<sup>th</sup> chunk (after 1 GB of uncompressed data).

Field	Comments	Size
E8 translation	0-disabled, 1-enabled	1 bit
Translation size high word	Only present if enabled	0 or 16 bits
Translation size low word	Only present if enabled	0 or 16 bits

## 2.12 Chunk Size

The LZXD compressor emits chunks of compressed data, each of which represents exactly 32 KB of uncompressed data until the last chunk in the stream, which can represent less than 32 KB. In order to ensure that an exact number of input bytes represent an exact number of output bytes for each chunk, after each 32 KB of uncompressed data is represented in the output compressed bitstream, the output bitstream is padded with up to 15 bits of zeros to realign the bitstream on a 16-bit boundary (even byte boundary) for the next 32 KB of data. This results in a compressed chunk of a byte-aligned size. The compressed chunk could be significantly smaller than 32 KB or possibly larger than 32 KB if the data is incompressible.

The LZXD engine encodes a byte-aligned little-endian 16-bit compressed chunk size prefix field preceding each compressed chunk in the compressed byte stream. The chunk prefix chain could be followed in the compressed stream without decompressing any data. The next

chunk prefix is at location computed by absolute byte offset location of this chunk prefix plus 2 (for the size of the chunk size prefix field) plus the current chunk size.

### 2.13 Block Header

An LZXD Block represents a sequence of compressed data that is encoded with the same set of Huffman trees, or a sequence of uncompressed data. There can be one or more LZXD Blocks in a compressed stream, each with its own set of Huffman trees. Blocks need not start or end on a chunk boundary; blocks can span multiple chunks, or a single chunk can contain multiple blocks. The Block Type field indicates which type of block follows, and the Block Size field indicates the number of uncompressed bytes represented by the block. Following the generic Block Header, there is a type-specific header describing the remainder of the block.

Field	Comments	Size
Block Type	See valid values below	3 bits
Block Size MSB	Block size high 8 bits of 24	8 bits
Block Size byte 2	Block size middle 8 bits of 24	8 bits
Block Size LSB	Block size low 8 bits of 24	8 bits

### 2.14 Block Type

Each block of compressed data begins with a 3-bit field indicating the block type, followed by the Block Size and then type-specific Block Data. Of the eight values possible, only three are valid types:

Bits	Value	Meaning
001	1	Verbatim block
010	2	Aligned offset block
011	3	Uncompressed block
other	0, 4-7	Invalid

### 2.15 Block Size

The Block Size field indicates the number of uncompressed bytes represented by the block. The maximum Block Size is  $2^{24}-1$  (16MB-1 or 0x00FFFFFF). The Block Size is encoded in the bitstream as three 8-bit fields comprising a 24-bit value, most significant to least significant, immediately following the Block Type encoding.

### 2.15.1 Uncompressed Block

Following the generic Block Header, an uncompressed block begins with 1 to 16 bits of zero padding to align the bit buffer on a 16-bit boundary. At this point, the bitstream ends, and a *byte stream* begins. Following the zero padding, new 32-bit values for R0, R1, and R2 are output in little-endian form, followed by the uncompressed data bytes themselves. Finally, if the uncompressed data length is odd, one extra byte of zero padding is encoded to re-align the following bitstream.

Field	Comments	Size
Padding to align following field on 16-bit boundary	Bits have value of zero	Variable, 1..16 bits

Then, the following fields are then encoded directly in the byte stream, NOT the bitstream of byte-swapped 16-bit words:

R0	LSB to MSB (little endian dword)	4 bytes
R1	LSB to MSB (little endian dword)	4 bytes
R2	LSB to MSB (little endian dword)	4 bytes
Uncompressed raw data bytes	Can use direct memcpy	1...2 <sup>24</sup> -1 bytes
Padding to re-align bitstream	Only if uncompressed size is odd	0 or 1 byte

Then the bitstream of byte-swapped 16 bit integers resumes for the next Block Type field (if there are subsequent blocks).

The decoded R0, R1, and R2 values are used as initial Repeated Offset values to decode the subsequent compressed block if present.

### 2.15.2 Verbatim Block

A verbatim block consists of the following fields following the generic block header:

Entry	Comments	Size
Pre-tree for first 256 elements of main tree	20 elements, 4 bits each	80 bits
Path lengths of first 256 elements of main tree	Encoded using pre-tree	Variable
Pre-tree for remainder of main tree	20 elements, 4 bits each	80 bits
Path lengths of remaining elements of main tree	Encoded using pre-tree	Variable
Pre-tree for length tree	20 elements, 4 bits each	80 bits
Path lengths of elements in length tree	Encoded using pre-tree	Variable
Token sequence (matches and literals)	Described later	Variable

### 2.15.3 Aligned Offset Block

An aligned offset block consists of the following, the only difference from Verbatim header being the existence of the Aligned Offset Tree preceding the other trees.

Entry	Comments	Size
Aligned offset tree	8 elements, 3 bits each	24 bits
Pre-tree for first 256 elements of main tree	20 elements, 4 bits each	80 bits
Path lengths of first 256 elements of main tree	Encoded using pre-tree	Variable
Pre-tree for remainder of main tree	20 elements, 4 bits each	80 bits
Path lengths of remaining elements of main tree	Encoded using pre-tree	Variable
Pre-tree for length tree	20 elements, 4 bits each	80 bits
Path lengths of elements in length tree	Encoded using pre-tree	Variable
Token sequence (matches and literals)	Described later	Variable

### 2.15.4 Encoding the Trees and Pre-Trees

Because all trees used in LZXD are created in the form of a canonical Huffman tree, the path length of each element in the tree is sufficient to reconstruct the original tree. The main tree and the length tree are each encoded using the method described here. However, the main tree is encoded in two components as if it were two separate trees, the first tree corresponding to the first 256 tree elements (uncompressed symbols), and the second tree corresponding to the remaining elements (matches).

Since trees are output several times during compression of large amounts of data (multiple blocks), LZX optimizes compression by encoding only the delta path lengths between the current and previous trees. In the case of the very first such tree, the delta is calculated against a tree in which all elements have a zero path length.

Each tree element can have a path length from 0 to 16 (inclusive) where a zero path length indicates that the element has a zero frequency and is not present in the tree. Tree elements are output in sequential order starting with the first element. Elements can be encoded in one of two ways: If several consecutive elements have the same path length, then run length encoding is employed; otherwise the element is output by encoding the difference between the current path length and the previous path length of the tree, mod 17. These output methods are described in the following table:

## Tree codes

Code	Operation
0-16	$\text{Len}[x] = (\text{prev\_len}[x] + \text{code}) \bmod 17$
17	Zeroes = getbits(4) $\text{Len}[x] = 0$ for next (4 + Zeroes) elements
18	Zeroes = getbits(5) $\text{Len}[x] = 0$ for next (20 + Zeroes) elements
19	Same = getbits(1) Decode new Code $\text{Value} = (\text{prev\_len}[x] + \text{Code}) \bmod 17$ $\text{Len}[x] = \text{Value}$ for next (4 + Same) elements

Each of the 17 possible values of  $(\text{len}[x] - \text{prev\_len}[x]) \bmod 17$ , plus three additional codes used for run-length encoding, are *not* output directly as 5 bit numbers, but are instead encoded via a Huffman tree called the *pre-tree*. The pre-tree is generated dynamically according to the frequencies of the 20 allowable tree codes. The structure of the pre-tree is encoded in a total of 80 bits by using 4 bits to output the path length of each of the 20 pre-tree elements. Once again, a zero path length indicates a zero frequency element.

### Pre-tree

Length of tree code 0	4 bits
Length of tree code 1	4 bits
Length of tree code 2	4 bits
...	...
Length of tree code 18	4 bits
Length of tree code 19	4 bits

The “real” tree is then encoded using the pre-tree Huffman codes.

### 2.15.5 Compressed Token Sequence

The compressed token sequence (bitstream) contains the Huffman-encoded matches and literals using the Huffman trees specified in the block header. Decompression continues until the number of decompressed bytes corresponds exactly to the number of uncompressed bytes indicated in the block header.

The representation of an unmatched literal character in the output is simply the appropriate element index 0...255 from the Main Huffman Tree.

The representation of a match in the output involves several transformations, as shown in the following diagram. At the top of the diagram are the match length (2..257) and the match offset (0...WINDOW\_SIZE-4). The match offset and match length are split into sub-components and encoded separately. For matches of length 257..32768, the token indicates

match length 257 and then there is an additional Extra Length value encoded in the bitstream following the other Match subcomponent fields.

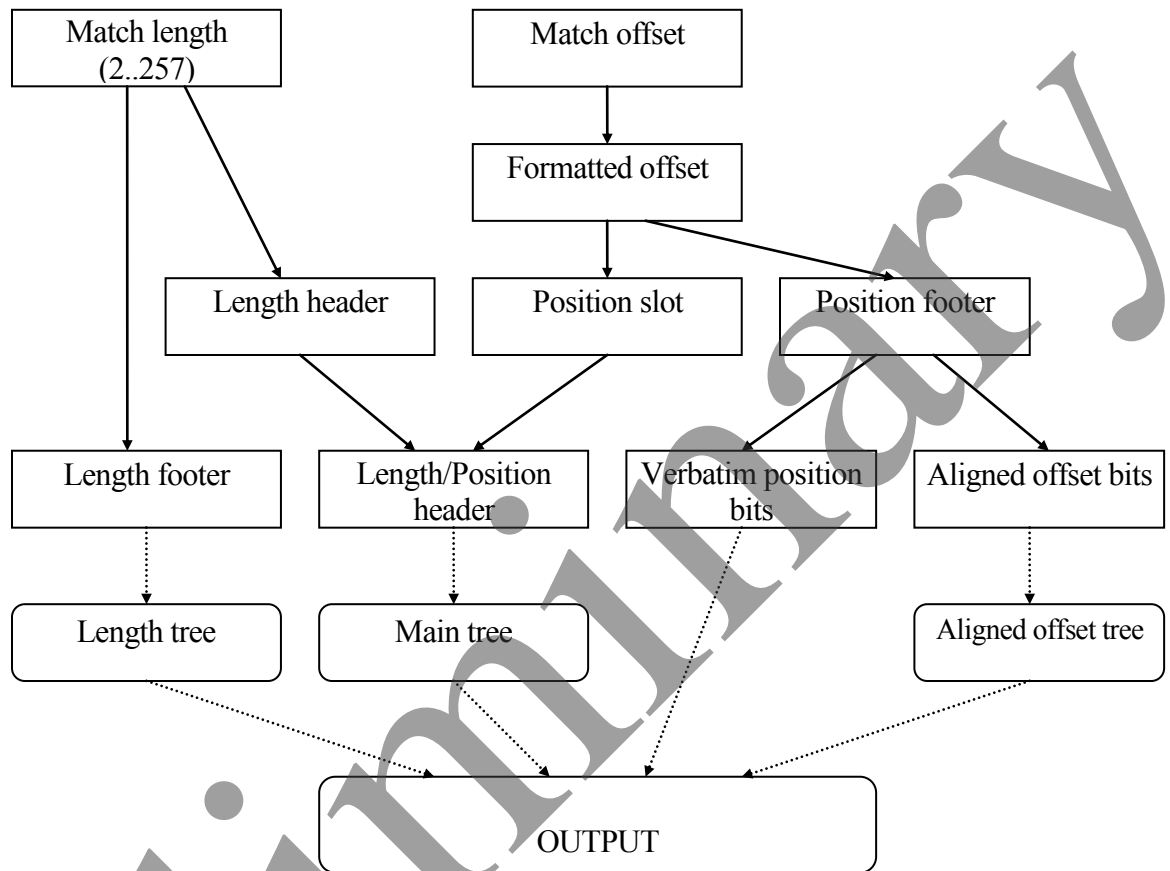


Figure 1: Diagram of match encoding sub components

### 2.15.6 Converting Match Offset into Formatted Offset Values

The match offset, range 1...(WINDOW\_SIZE-4), is converted into a *formatted offset* by determining whether the offset can be encoded as a repeated offset, as shown in the following pseudocode. It is acceptable to not encode a match as a repeated offset even if it is possible to do so.

#### Converting a match offset to a formatted offset

```

if offset == R0 then
    formatted offset ← 0
else if offset == R1 then
    formatted offset ← 1
else if offset == R2 then

```

```

        formatted offset ← 2
    else
        formatted offset ← offset + 2
    endif

```

### 2.15.7 Converting Formatted Offset into Position Slot and Position Footer Values

The formatted offset is subdivided into a position slot and position footer. The position slot defines the most significant bits of the formatted offset in the form of a base position as shown in the table on the following page. The position footer defines the remaining least significant bits of the formatted offset. As the following table shows, the number of bits dedicated to the position footer grows as the formatted offset becomes larger, meaning that each position slot addresses a larger and larger range.

The number of position slots available depends on the window size. The number of bits of position footer for each position slot is fixed and also shown in the table below.

**Position Slot Table (Formatted Offset = Base Position of Slot + Footer Bits Value)**

Position slot number	Base position	Footer bits	Base plus position footer range
0 (R0)	0	0	0
1 (R1)	1	0	1
2 (R2)	2	0	2
3 (offset 1)	3	0	3
4 (offset 2..3)	4	1	4-5
5 (offset 4..5)	6	1	6-7
6 (offset 6..9)	8	2	8-11
7 (..etc..)	12	2	12-15
8	16	3	16-23
9	24	3	24-31
10	32	4	32-47
11	48	4	48-63
12	64	5	64-95
13	96	5	96-127
14	128	6	128-191
15	192	6	192-255
16	256	7	256-383
17	384	7	384-511
18	512	8	512-767
19	768	8	768-1023
20	1024	9	1024-1535
21	1536	9	1536-2047



22	2048	10	2048-3071
23	3072	10	3072-4095
24	4096	11	4096-6143
25	6144	11	6144-8191
26	8192	12	8192-12287
27	12288	12	12288-16383
28	16384	13	16384-24575
29	24576	13	24576-32767
30	32768	14	32768-49151
31	49152	14	49152-65535
32	65536	15	65536-98303
33	98304	15	98304-131071
34	131072	16	131072-196607
35	196608	16	196608-262143
36	262144	17	262144-393215
37	393216	17	393216-524287
38	524288	17	524288-655359
39	655360	17	655360-786431
40	786432	17	786432-917503
41	917504	17	917504-1048575
42	1048576	17	1048576-1179647
..etc..	..etc..	17 (all)	..etc..
288	33292288	17	33292288-33423359
289	33423360	17	33423360-33554431

### 2.15.8 Converting Position Footer into Verbatim Bits or Aligned Offset Bits

The position footer can be further subdivided into verbatim bits and aligned offset bits if the current block type is “aligned offset”. If the current block is not an aligned offset block then there are no aligned offset bits, and the verbatim bits are the position footer.

If aligned offsets are used, then the lower 3 bits of the position footer are the aligned offset bits, while the remaining portion of the position footer are the verbatim bits. In the case where there are less than 3 bits in the position footer (for example, formatted offset is  $\leq 15$ ) it is not possible to take the “lower 3 bits of the position footer” and therefore there are no aligned offset bits, and the verbatim bits and the position footer are the same.

Pseudocode for splitting position footer into verbatim bits and aligned offset

```

if block_type is aligned_offset_block then
  if formatted_offset <= 15 then
    verbatim_bits ← position_footer

```

```

        aligned_offset ← null
    else
        aligned_offset ← position_footer
        verbatim_bits ← position_footer >> 3
    endif
else
    verbatim_bits ← position_footer
    aligned_offset ← null
endif

```

### 2.15.9 Converting Match Length into Length Header and Length Footer Values

The match length is converted into a length header and a length footer. The length header can have one of eight possible values, from 0...7 (inclusive), indicating a match of length 2, 3, 4, 5, 6, 7, 8, or a length greater than 8. If the match length is 8 or less, then there is no length footer. Otherwise, the value of the length footer is equal to the match length minus 9.

Pseudocode for obtaining the length header and footer

```

if match_length <= 8
    length_header ← match_length-2
    length_footer ← null
else
    length_header ← 7
    length_footer ← match_length-9
endif

```

Example conversions of some match lengths to header and footer values

Match length	Length header	Length footer value
2	0	None
3	1	None
4	2	None
5	3	None
6	4	None
7	5	None
8	6	None
9	7	0
10	7	1
50	7	41
256	7	247
257 or larger	7	248

## 2.15.10 Converting Length Header and Position Slot into Length/Position Header Values

The Length/Position header is the stage which correlates the match position with the match length (using only the most significant bits), and is created by combining the length header and the position slot as shown below:

$$\text{len\_pos\_header} \leftarrow (\text{position\_slot} \ll 3) + \text{length\_header}$$

This operation creates a unique value for every combination of match length 2, 3, 4, 5, 6, 7, 8 with every possible position slot. The remaining match lengths greater than 8 are all lumped together, and as a group are correlated with every possible position slot.

## 2.16 Extra Length

If the match length is 257 or larger, the encoded match length token value is 257, and an encoded Extra Length field follows the other match encoding components in the bitstream.

### Extra Length Encoding

Prefix	Number of Bits to Decode	Base Value to Add to Decoded Value
0	8	257
10	10	257 + 256
110	12	257 + 256 + 1024
111	15	257

### 2.16.1 Encoding a Match

The match is finally output in up to five components, in the following order:

1. Main Tree element at index ( $\text{len\_pos\_header} + 256$ ).
2. If  $\text{length\_footer} \neq \text{null}$ , then Length Tree element  $\text{length\_footer}$ .
3. If  $\text{verbatim\_bits} \neq \text{null}$ , then output  $\text{verbatim\_bits}$ .
4. If  $\text{aligned\_offset\_bits} \neq \text{null}$ , then output element  $\text{aligned\_offset}$  from the aligned offset tree.
5. If match length 257 or larger, output appropriate Extra Length prefix and value.

## 2.17 Encoding a Literal

A literal byte not part of a match is encoded simply as a Main Tree element index 0..256 corresponding to the value of the literal byte.

### 2.17.1 Decoding Matches and Literals (Aligned and Verbatim Blocks)

Decoding is performed by first decoding an element from the Main Tree and then, if the item is a match, determining which additional components are necessary to decode to reconstruct the match. Pseudocode for decoding a match or an uncompressed character is shown here:

```
main_element = main_tree.decode_element()

if (main_element < 256 ) /* is a literal character */

    window[ curpos ] ← (byte) main_element
    curpos ← curpos + 1

else /* is a match */

    length_header ← (main_element - 256) & 7

    if (length_header == 7)
        match_length ← length_tree.decode_element() + 7 + 2
    else
        match_length ← length_header + 2 /* no length footer */
    endif

    position_slot ← (main_element - 256) >> 3

    /* check for repeated offsets (positions 0,1,2) */
    if (position_slot == 0)
        match_offset ← R0
    else if (position_slot == 1)
        match_offset ← R1
        swap(R0 ↔ R1)
    else if (position_slot == 2)
        match_offset ← R2
        swap(R0 ↔ R2)
    else /* not a repeated offset */
        offset_bits ← footer_bits[ position_slot ]

        if (block_type == aligned_offset_block)
            if (offset_bits >= 3) /* this means there are some aligned
                bits */
                verbatim_bits ← (readbits(offset_bits-3)) << 3
                aligned_bits ← aligned_offset_tree.decode_element();
            else /* 0, 1, or 2 verbatim bits */
                verbatim_bits ← readbits(offset_bits)
                aligned_bits ← 0
            endif

            formatted_offset ← base_position[ position_slot ]
                + verbatim_bits + aligned_bits

        else /* block_type == verbatim_block */
```

```

        verbatim_bits ← readbits(offset_bits)
        formatted_offset ← base_position[ position_slot ] +
            verbatim_bits
    endif

    match_offset ← formatted_offset - 2

    /* update repeated offset LRU queue */
    R2 ← R1
    R1 ← R0
    R0 ← match_offset

endif

/* check for extra length */

if (match_length == 257)
    if (readbits( 1 ) != 0)
        if (readbits( 1 ) != 0)
            if (readbits( 1 ) != 0)
                extra_len = readbits( 15 )
            else
                extra_len = readbits( 12 ) + 1024 + 256
            endif
        else
            extra_len = readbits( 10 ) + 256
        endif
    else
        extra_len = readbits( 8 )
    endif

    match_length ← 257 + extra_len

endif

/* copy match data */
for (i = 0; i < match_length; i++)
    window[curpos + i] ← window[curpos + i - match_offset]

curpos ← curpos + match_length

endif

```

### 3 Protocol Examples

Below is an example of a sample encoding sequence of a simple 3-byte text input “abc” encoded as uncompressed block type.

Bits to Decode	Value of Decoded Bits	Interpretation
16	0x0014	Chunk Size: 20 bytes
1	0	E8 Translation:disabled

3	3 (binary 011)	Block Type: uncompressed
24	0x000003	Block Size: 3 bytes
4	binary 0000	Padding to word-align following
4 bytes	0x00000001 (little-endian dword)	R0: 1
4 bytes	0x00000001 (little-endian dword)	R1: 1
4 bytes	0x00000001 (little-endian dword)	R2: 1
4 bytes	0x61, 0x62, 0x63	Uncompressed bytes: "abc"
1 byte	0x00	Padding to restore word-alignment

This is the raw hexadecimal compressed byte sequence of the above encoded fields:

14 00 00 30 30 00 01 00 00 00 01 00 00 00 01 00 00 00 61 62 63 00

#### 4 Appendix A: Office/Exchange Behavior

The information in this specification is applicable to the following versions of Office/Exchange:

- Office 2003 with Service Pack 3 applied
- Exchange 2003 with Service Pack 2 applied
- Office 2007 with Service Pack 1 applied
- Exchange 2007 with Service Pack 1 applied

Exceptions, if any, are noted below. Unless otherwise specified, any statement of optional behavior in this specification prescribed using the terms SHOULD or SHOULD NOT implies Office/Exchange behavior in accordance with the SHOULD or SHOULD NOT prescription. Unless otherwise specified, the term MAY implies Office/Exchange does not follow the prescription.

## Index

- Bitstream, 5
- Block header, 11
- Block size, 11
- Block type, 11
- Chunk size, 10
- Description, 4
  - Bitstream, 5
  - Block header, 11
  - Block size, 11
  - Block type, 11
  - Chunk size, 10
  - E8 call translation, 8
  - Encoding a literal, 19
  - Extra length, 19
  - Huffman trees, 6
  - LZ77, 4
  - LZX, 5
  - LZXD, 5
  - Match lengths, 8
  - Position slot, 7
  - Reference data, 5
  - Repeated offsets, 7
  - Window size, 5
- E8 call translation, 8
- Encoding a literal, 19
- Extra length, 19
- Glossary, 4
- Huffman trees, 6
- Informative references, 4
- LZ77, 4
- LZX, 5
- LZXD, 5
- Match lengths, 8
- Normative references, 4
- Office/Exchange behavior, 22
- Position slot, 7
- Protocol examples, 21
- Reference data, 5
- References, 4
  - Informative references, 4
  - Normative references, 4

Repeated offsets, 7  
Window size, 5

Preliminary